

Inverses of Multivariable Polynomial Matrices by Discrete Fourier Transforms

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Outline of the presentation

- nD polynomial interpolation
nD DFT
Polynomial interpolation and DFT
- Computation of the Generalized inverse via DFT
- Computation of the Drazin inverse via DFT
- Complexity
Perturbation analysis
- Conclusions



Objectives

- A **DFT** based algorithm for the evaluation of the **generalized inverse** and the **Drazin inverse** of a multivariable polynomial matrix is presented.
- The **efficiency** of the algorithms is illustrated via complexity analysis



Notation

Consider the polynomial matrix with real coefficients in the n indeterminates z_1, z_2, \dots, z_n (called nD polynomial matrix)

$$A(z_1, \dots, z_n) = \sum_{k_1=0}^{M_1} \sum_{k_2=0}^{M_2} \dots \sum_{k_n=0}^{M_n} \left(A_{k_1 \dots k_n} \right) \times \left(z_1^{k_1} \dots z_n^{k_n} \right) \in R[z_1, \dots, z_n]^{p \times m}$$

with $A_{k_1 \dots k_n} \in R^{p \times m}$. For brevity and when the variables are obvious, we will use the notation \bar{z} instead of (z_1, z_2, \dots, z_n) .



nD polynomial interpolation

Assume that we need to compute the value of $f(A)$ where f is some polynomial function. The evaluation-interpolation method of computing $f(A)$ is based on the following steps

- **Step 1** Evaluation of the polynomial (matrix) at a set of $R = \prod_{i=1}^n (M_i + 1)$ suitably chosen points. This step results in a set of constant matrices.
- **Step 2** Application of the function f on the set of constant matrices derived from step 1.
- **Step 3** Computation of the coefficients of $f(A)$ through interpolation.



nD DFT

Consider the finite sequence $X(k_1, \dots, k_n)$ and $\tilde{X}(r_1, \dots, r_n)$, $k_i, r_i = 0, 1, \dots, M_i$. In order for the sequence $X(k_1, \dots, k_n)$ and $\tilde{X}(r_1, \dots, r_n)$ to constitute a DFT pair the following relations should hold [Dudgeon] :

$$\tilde{X}(r_1, \dots, r_n) = \sum_{k_1=0}^{M_1} \sum_{k_2=0}^{M_2} \dots \sum_{k_n=0}^{M_n} X(k_1, \dots, k_n) W_1^{-k_1 r_1} \dots W_1^{-k_n r_n}$$

$$X(k_1, \dots, k_n) = \frac{1}{R} \sum_{r_1=0}^{M_1} \sum_{r_2=0}^{M_2} \dots \sum_{r_n=0}^{M_n} \tilde{X}(r_1, \dots, r_n) W_1^{k_1 r_1} \dots W_1^{k_n r_n}$$

where

$$W_i = e^{\frac{2\pi j}{M_i+1}} \quad \forall i = 1, 2, 3, \dots, n$$

$$R = \prod_{i=1}^n (M_i + 1)$$

and X, \tilde{X} are discrete argument matrix-valued functions, with dimensions $p \times m$.



Polynomial interpolation and DFT

Efficiency of steps 1, 2 and 3 of the evaluation-interpolation algorithm.

- In step 2 a suitable algorithm computing $f(A)$ where A is a constant matrix, must be chosen. Then, the critical points of the algorithm are steps 1 and 3.
- By using as evaluation points in **Step 1** *Fourier points*, the evaluation of the polynomial matrix is equivalent to the **DFT** of the multidimensional matrix of the coefficients. **Step 3** becomes an **inverse DFT** problem.
- By using FFT the speed of the algorithm improves dramatically.



Generalized Inverse of a Multivariable Polynomial Matrix

[Penrose] For every matrix $A \in R^{p \times m}$, a unique matrix $A^+ \in R^{m \times p}$, which is called generalized inverse, exists satisfying

(i) $AA^+A = A$

(ii) $A^+AA^+ = A^+$

(iii) $(AA^+)^T = AA^+$

(iv) $(A^+A)^T = A^+A$

where A^T denotes the transpose of A . In an analogous way we define the generalized inverse $A(z_1, \dots, z_n)^+ \in R(z_1, \dots, z_n)^{m \times p}$ of a multivariable polynomial matrix

$$A(z_1, \dots, z_n) \in R[z_1, \dots, z_n]^{p \times m}$$

Generalized Inverse of a Multivariable Polynomial Matrix

Following the steps of [Karampetakis 1997] we have:

Let $A(\bar{z}) = A(z_1, \dots, z_n) \in R[z_1, \dots, z_n]^{p \times m}$ and

$$\begin{aligned} a(s, z_1, \dots, z_n) &= \det \left[sI_p - A(\bar{z})A(\bar{z})^T \right] \\ &= \left(a_0(\bar{z})s^p + \dots + a_{p-1}(\bar{z})s + a_p(\bar{z}) \right) \end{aligned}$$

$a_0(\bar{z}) = 1$, be the characteristic polynomial of $A(\bar{z}) \times A(\bar{z})^T$.

Let k such that it satisfies $a_p(\bar{z}) \equiv 0, \dots, a_{k+1}(\bar{z}) \equiv 0$ while $a_k(\bar{z}) \neq 0$, and define

$$\Lambda := \{(\bar{z}) \in C^n : a_k(\bar{z}) = 0\}.$$

Then the generalized inverse $A(\bar{z})^+$ of $A(\bar{z})$ for $\bar{z} \in C^n - \Lambda$ is given by

$$A(\bar{z})^+ = -\frac{1}{a_k(\bar{z})} A(\bar{z})^T B_{k-1}(\bar{z})$$

$$B_{k-1}(\bar{z}) = a_0(\bar{z}) \left[A(\bar{z})A(\bar{z})^T \right]^{k-1} + \dots + a_{k-1}(\bar{z}) I_p$$

Generalized Inverse of a Multivariable Polynomial Matrix

Step 1

$$a(s, z_1, \dots, z_n) = \sum_{k_0=0}^{b_0} \sum_{k_1=0}^{b_1} \dots \sum_{k_n=0}^{b_n} \left(a_{k_0 k_1 \dots k_n} \right) \left(s^{k_0} z_1^{k_1} \dots z_n^{k_n} \right)$$

- Define

$$\deg_s \left(a(s, z_1, \dots, z_n) \right) = p := b_0$$

$$\deg_{z_1} \left(a(s, z_1, \dots, z_n) \right) \leq 2pM_1 := b_1$$

⋮

$$\deg_{z_n} \left(a(s, z_1, \dots, z_n) \right) \leq 2pM_n := b_n$$

- We use following $R_1 = \prod_{i=0}^n (b_i + 1)$ interpolation points

$$u_i(r_j) = W_i^{-r_j}; i = 0, \dots, n \text{ and } r_j = 0, 1, \dots, b_i \text{ where } W_i = e^{\frac{2\pi j}{b_i+1}}$$

$$\tilde{a}_{r_0 r_1 \dots r_n} = \det \left[u_0(r_0) I_p - A(u_1(r_1), \dots, u_n(r_n)) \left[A(u_1(r_1), \dots, u_n(r_n)) \right]^T \right] = \sum_{l_0=0}^{b_0} \sum_{l_1=0}^{b_1} \dots \sum_{l_n=0}^{b_n} \left(a_{l_0 l_1 \dots l_n} \right) \left(W_0^{-r_0 l_0} \dots W_n^{-r_n l_n} \right)$$

- $[a_{l_0 l_1 \dots l_n}]$ and $[\tilde{a}_{r_0 r_1 \dots r_n}]$ form a DFT pair Thus we compute $a_{l_0 l_1 \dots l_n}$ via IDFT



Generalized Inverse of a Multivariable Polynomial Matrix Step 2

Find using a loop k such that $a_{k+1}(\bar{z}) = a_{k+2}(\bar{z}) = \cdots = a_p(\bar{z}) = 0$ and $a_k(\bar{z}) \neq 0$

Generalized Inverse of a Multivariable Polynomial Matrix Step 3

$$C(\bar{z}) = A(\bar{z})^T B_{k-1}(\bar{z}) = a_0(\bar{z}) \left[A(\bar{z}) A(\bar{z})^T \right]^{k-1} + \dots + a_{k-1}(\bar{z}) I_p$$

- $$C(\bar{z}) = \sum_{l_1=0}^{n_1} \dots \sum_{l_n=0}^{n_n} C_{l_0 \dots l_n} \left(z_1^{l_1} \dots z_n^{l_n} \right)$$

where $n_i = \max \{ 2(k-1)M_i + M_i, k = 1, \dots, p \} = (2p-1)M_i$

- We use the following $R_2 = \prod_{i=1}^n \{ (2p-1)M_i + 1 \}$ points

$$u_i(r_j) = W_i^{-r_j}; i = 1, \dots, n \text{ and } r_j = 0, 1, \dots, n_i \text{ where } W_i = e^{\frac{2\pi j}{n_i+1}}$$

$$\tilde{C}_{r_1 \dots r_n} = \sum_{l_1=0}^{n_1} \dots \sum_{l_n=0}^{n_n} C_{l_0 \dots l_n} W_1^{-r_1 l_1} \dots W_n^{-r_n l_n}$$

- $\left[C_{l_0 \dots l_n} \right]$ and $\left[\tilde{C}_{r_1 \dots r_n} \right]$ form a DFT pair. Thus we compute $C_{l_0 \dots l_n}$ via IDFT



Generalized Inverse of a Multivariable Polynomial Matrix Step 4

$$A(\bar{z})^+ = -\frac{1}{a_k(\bar{z})}C(\bar{z})$$



Drazin inverse

For every matrix $A \in R^{m \times m}$, there exists a unique matrix $A^D \in R^{m \times m}$, which is called Drazin inverse, satisfying

(i) $A^D A^{k+1} = A^k$ for $k = \text{ind}(A) = \min(k \in N : \text{rank}(A^k) = \text{rank}(A^{k+1}))$

(ii) $A^D A A^D = A^D$

(iii) $A A^D = A^D A$



Drazin inverse

Consider a nonregular $n \times n$ polynomial matrix $A(\bar{z})$. Assume that

$$a(s, z_1, \dots, z_n) = \det [sI_m - A(\bar{z})] = (a_0(\bar{z})s^m + \dots + a_{m-1}(\bar{z})s + a_m(\bar{z}))$$

where $a_0(\bar{z}) \equiv 1, z \in C$ is the characteristic polynomial of $A(\bar{z})$.

Also, consider the following sequence of $m \times m$ polynomial matrices

$$B_j(\bar{z}) = a_0(\bar{z})A(\bar{z})^j + \dots + a_{j-1}(\bar{z})A(\bar{z}) + a_j(\bar{z})I_m, a_0(\bar{z}) = 1, j = 0, \dots, m$$

Let $a_m(\bar{z}) \equiv 0, \dots, a_{t+1}(\bar{z}) \equiv 0, a_t(\bar{z}) \neq 0$.

Define the following set $\Lambda = \{\bar{z}_i \in C^n : a_t(\bar{z}_i) = 0\}$

Also, assume that $B_m(\bar{z}), \dots, B_r(\bar{z}) = 0, B_{r-1}(\bar{z}) \neq 0$ and $k = r - t$.

In the case $\bar{z} \in C^n - \Lambda$ and $k > 0$, the Drazin inverse of $A(\bar{z})$ is given by

$$A(\bar{z})^D = \frac{A(\bar{z})^k B_{t-1}(\bar{z})^{k+1}}{a_t(\bar{z})^{k+1}}$$

$$B_{t-1}(\bar{z}) = a_0(\bar{z})A(\bar{z})^{t-1} + \dots + a_{t-2}(\bar{z})A(\bar{z}) + a_{t-1}(\bar{z})I_m$$

In the case $\bar{z} \in C^n - \Lambda$ and $k = 0$, we get $A(\bar{z})^D = O$.

For $\bar{z}_i \in \Lambda$ we can use the same algorithm again.

Drazin Inverse of a Multivariable Polynomial Matrix

Step 1

$$a(s, z_1, \dots, z_n) = \sum_{k_0=0}^{b_0} \sum_{k_1=0}^{b_1} \dots \sum_{k_n=0}^{b_n} \left(a_{k_0 k_1 \dots k_n} \right) \left(s^{k_0} z_1^{k_1} \dots z_n^{k_n} \right)$$

- Define

$$\deg_s \left(a(s, z_1, \dots, z_n) \right) = m := b_0$$

$$\deg_{z_1} \left(a(s, z_1, \dots, z_n) \right) \leq m M_1 := b_1$$

⋮

$$\deg_{z_n} \left(a(s, z_1, \dots, z_n) \right) \leq m M_n := b_n$$

- We use following $R_1 = \prod_{i=0}^n (b_i + 1)$ interpolation points

$$u_i(r_j) = W_i^{-r_j}; i = 0, \dots, n \text{ and } r_j = 0, 1, \dots, b_i \text{ where } W_i = e^{\frac{2\pi j}{b_i+1}}$$

$$\tilde{a}_{r_0 r_1 \dots r_n} = \det \left[u_0(r_0) I_p - A(u_1(r_1), \dots, u_n(r_n)) \right]$$

- $[a_{l_0 l_1 \dots l_n}]$ and $[\tilde{a}_{r_0 r_1 \dots r_n}]$ form a DFT Thus we compute $a_{l_0 l_1 \dots l_n}$ via IDFT



Drazin Inverse of a Multivariable Polynomial Matrix Step 2

Find using a loop t such that $a_{t+1}(\bar{z}) = a_{t+2}(\bar{z}) = \dots = a_m(\bar{z}) = 0$ and $a_t(\bar{z}) \neq 0$

A polynomial matrix

$$B(z_1, \dots, z_n) \in R[z_1, \dots, z_n]^{m \times m}$$

of degrees q_i in respect with variables $z_i, i = 1, \dots, n$, is the zero polynomial matrix iff its values at $R = \prod_{i=0}^n (q_i + 1)$ distinct points are the zero matrix.

Find r such that $B_m(\bar{z}), \dots, B_r(\bar{z}) = 0, B_{r-1}(\bar{z}) \neq 0$ and $k = r - t$ using the above lemma.

Drazin Inverse of a Multivariable Polynomial Matrix

Step 3

$$C(\bar{z}) = A(\bar{z})^k B_{t-1}(\bar{z})^{k+1}$$

- $C(\bar{z}) = \sum_{l_1=0}^{n_1} \dots \sum_{l_n=0}^{n_n} C_{l_0 \dots l_n} \left(z_1^{l_1} \dots z_n^{l_n} \right)$, $n_i = (t-1)(k+1)M_i$
- We use the following $R_2 = \prod_{i=1}^n \{(n_i + 1)\}$ points

$$u_i(r_j) = W_i^{-r_j}; i = 1, \dots, n \text{ and } r_j = 0, 1, \dots, n_i \text{ where } W_i = e^{\frac{2\pi j}{n_i+1}}$$

$$\tilde{C}_{r_1 \dots r_n} = \sum_{l_1=0}^{n_1} \dots \sum_{l_n=0}^{n_n} C_{l_0 \dots l_n} W_1^{-r_1 l_1} \dots W_n^{-r_n l_n}$$

- $\left[C_{l_0 \dots l_n} \right]$ and $\left[\tilde{C}_{r_1 \dots r_n} \right]$ form a DFT pair. Thus we compute $C_{l_0 \dots l_n}$ via IDFT

Drazin Inverse of a Multivariable Polynomial Matrix

Step 4

$$c(\bar{z}) = a_t(\bar{z})^{k+1}$$

- $c(\bar{z}) = \sum_{k_1=0}^{d_1} \cdots \sum_{k_n=0}^{d_n} (c_{k_1 \dots k_n}) (z_1^{k_1} \dots z_n^{k_n})$, where $d_i = tM_i(k+1)$
- We use the following $R_3 = \prod_{i=1}^n (d_i + 1)$ points

$$u_i(r_j) = W_i^{-r_j}; i = 1, \dots, n \text{ and } r_j = 0, 1, \dots, d_i, \text{ where } W_i = e^{\frac{2\pi j}{d_i+1}}$$

$$\tilde{c}_{r_1 \dots r_n} = \sum_{l_1=0}^{d_1} \cdots \sum_{l_n=0}^{d_n} (c_{k_1 \dots k_n}) (W_1^{-r_1 l_1} \dots W_n^{-r_n l_n})$$

- $[c_{l_0 \dots l_n}]$ and $[\tilde{c}_{r_1 \dots r_n}]$ form a DFT pair. Thus we compute $c_{l_0 \dots l_n}$ via IDFT



Drazin Inverse of a Multivariable Polynomial Matrix Step 5

$$A(\bar{z})^D = \frac{C(\bar{z})}{c(\bar{z})}$$



Complexity

- **Generalized inverse**

$O(mp^3 R_1 L_1)$ where

$$L_1 = (p+1) + \sum_{i=1}^n \log(2pM_i + 1) \text{ and } R_1 = (p+1) \prod_{i=1}^n (2pM_i + 1)$$



Complexity

- **Drazin inverse**

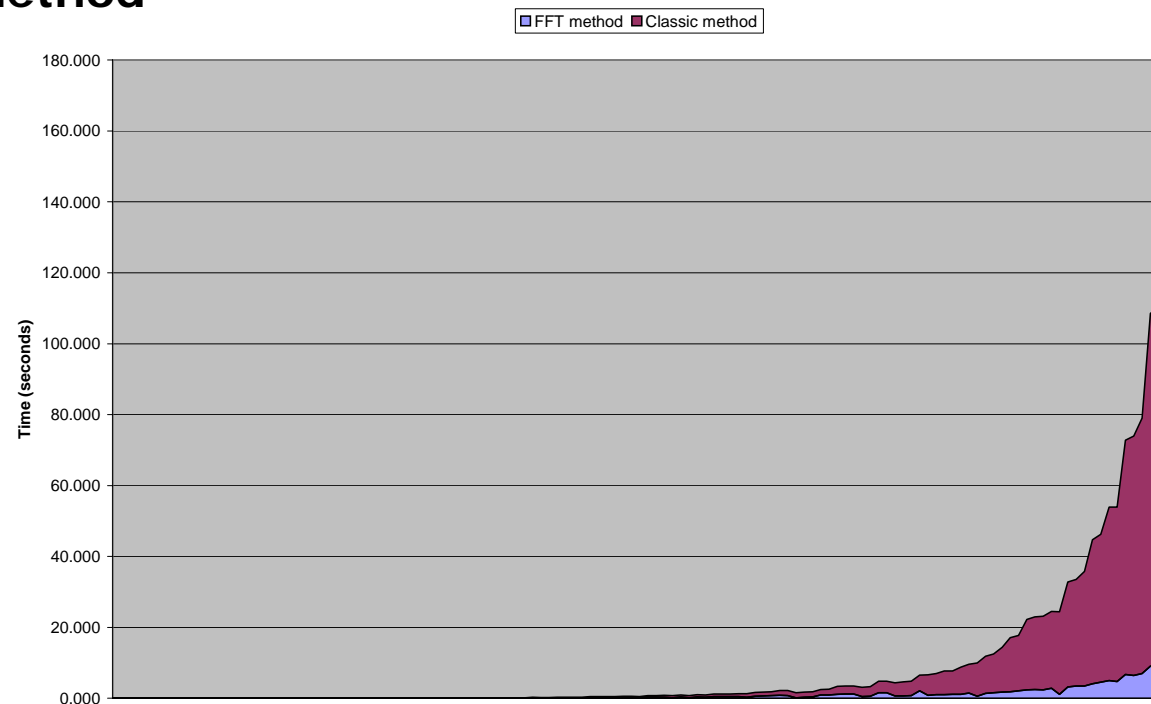
$O(m^4 RL)$ where

$$R = \max \{R_1, R_2, R_3\}, L = \max \{L_1, L_2, L_3\}$$

$$\text{and } L_1 = \sum_{i=0}^n \log(b_i + 1), L_2 = \sum_{i=1}^n \log(n_i + 1), L_3 = \sum_{i=1}^n \log(d_i + 1)$$

Numerical aspects of the algorithm

Comparison in Mathematica of the FFT and the symbolic method



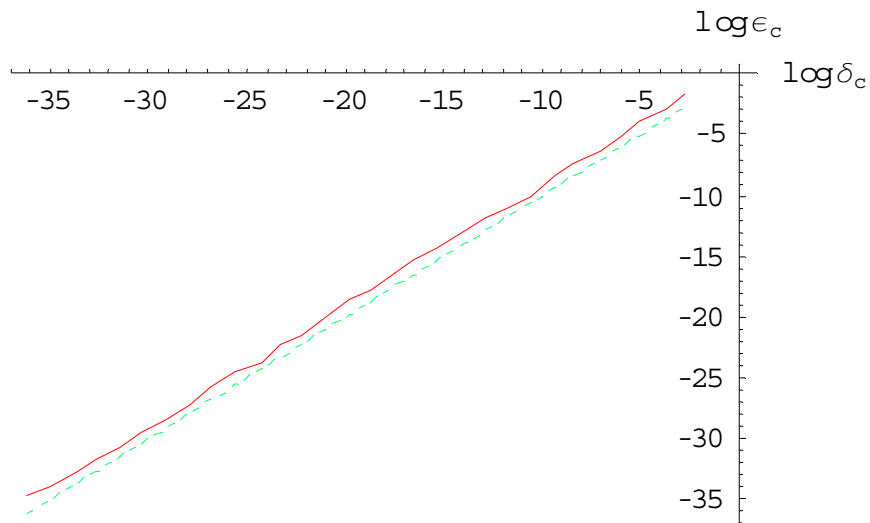
Random two variable polynomial matrices of dimensions up to **5x5** and degrees up to **8**.

Perturbation analysis

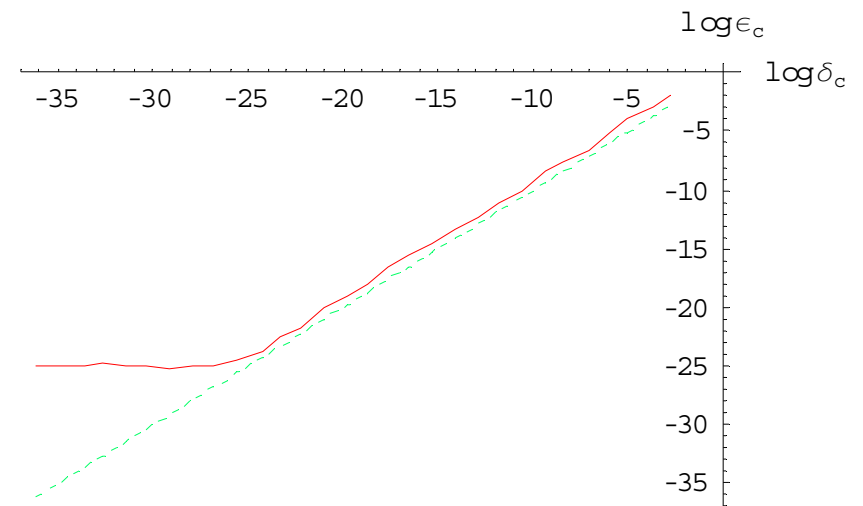
Relationship between the relative error $\varepsilon_c = \frac{\|(A(\bar{z}) + \delta_A(\bar{z}, c))^+ - A^+(\bar{z})\|_2}{\|A^+(\bar{z})\|_2}$ and the relative error in

the data $\delta_c = \frac{\|\delta_A(\bar{z}, c)\|_2}{\|A(\bar{z})\|_2}$, where $\|A(\bar{z})\|_2$ denotes the 2-norm of the coefficient matrix

$\begin{bmatrix} A_{0\dots 0} & \cdots & A_{M_1\dots M_n} \end{bmatrix}$ of $A(\bar{z})$.



Perturbation analysis of the denominator polynomial of the generalized inverse



Perturbation analysis of the numerator of the generalized inverse



Conclusions

- In this paper two algorithms have been presented for determining the Moore-Penrose and Drazin inverse of nD polynomial matrices.
- The algorithms are based on the fast Fourier transform and therefore have the main advantages of speed in contrast to other known algorithms.
- Applications include model matching, the solution of multivariable Diophantine equations and its application to control system synthesis problems, etc.
- The above mentioned algorithms may be easily extended in order to determine other kind of inverses such as $\{2\}$, $\{1,2\}$, $\{1,2,3\}$ and $\{1,2,4\}$ inverses of multivariable polynomial matrices.